

IS-ALGEBRAS WITH AN ADDITIONAL OPERATION

by

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ABSTRACT. In this paper we consider the algebraic counterpart of the fragment of the three-valued Lukasiewicz propositional calculus where the connectives are $\rightarrow\rightarrow$ (Lukasiewicz implication) and \wedge (conjunction). Then, we define a new equational class of algebras of type $(2,2,0)$, called IS-algebras with infimum (or SIS-algebras). They are an extention of the implicative three-valued Lukasiewicz algebras (or IS-algebras).

1. INTRODUCTION.

A.Rose [13] (see also [7]) gave a formalization of the implicative three-valued Lukasiewicz propositional calculus (or three-ILPC) by means of the following axiom schemes

$$(C1) \quad x \rightarrow\rightarrow (y \rightarrow\rightarrow x) ,$$

$$(C2) \quad (x \rightarrow\rightarrow y) \rightarrow\rightarrow ((y \rightarrow\rightarrow z) \rightarrow\rightarrow (x \rightarrow\rightarrow z)) ,$$

$$(C3) \quad ((x \rightarrow\rightarrow y) \rightarrow\rightarrow y) \rightarrow\rightarrow ((y \rightarrow\rightarrow x) \rightarrow\rightarrow x) ,$$

$$(C4) \quad ((x \rightarrow\rightarrow y) \rightarrow\rightarrow (y \rightarrow\rightarrow x)) \rightarrow\rightarrow (y \rightarrow\rightarrow x) ,$$

$$(C5) \quad ((x \rightarrow\rightarrow (x \rightarrow\rightarrow y) \rightarrow\rightarrow x) \rightarrow\rightarrow x) ,$$

and the rules of procedure

$$(R0) \quad \text{substitution rule} ,$$

$$(R1) \quad \text{"Modus ponens" rule} \quad \frac{x, x \rightarrow\rightarrow y}{y} .$$

Adopting the notations

$$(N1) \quad x \vee y = (x \rightarrow\rightarrow y) \rightarrow\rightarrow y ,$$

$$(N2) \quad x \rightarrow y = x \rightarrow\rightarrow (x \rightarrow\rightarrow y) .$$

we can write C3, C4 and C5 as follows

$$(C3) \quad (x \vee y) \rightarrow (y \vee x) ,$$

$$(C4) \quad (x \rightarrow y) \vee (y \rightarrow x) ,$$

$$(C5) \quad (x \rightarrow y) \vee x .$$

The followings formulas and rules are consequence of R0, R1 ,
C1,...,C5 (See [10]) .

$$(R2) \quad \frac{x}{x \rightarrow t} ,$$

$$(R3) \quad \frac{x \rightarrow y}{(y \rightarrow z) \rightarrow (x \rightarrow z)} ,$$

$$(C6) \quad x \rightarrow x ,$$

$$(C7) \quad (x \rightarrow y) \rightarrow ((z \rightarrow x) \rightarrow (z \rightarrow y)) ,$$

$$(R4) \quad \frac{(x \rightarrow y)}{(z \rightarrow x) \rightarrow (z \rightarrow y)} ,$$

$$(C8) \quad (x \rightarrow (x \rightarrow y)) \rightarrow (x \rightarrow y) ,$$

$$(C9) \quad (x \rightarrow (y \rightarrow z)) \rightarrow (y \rightarrow (x \rightarrow z)) ,$$

$$(R5) \quad \frac{(x \rightarrow (y \rightarrow z))}{(y \rightarrow (x \rightarrow z))} ,$$

$$(C10) \quad (x \rightarrow (y \rightarrow z)) \rightarrow ((x \rightarrow y) \rightarrow (x \rightarrow z)) ,$$

$$(C11) \quad x \rightarrow (y \rightarrow x) ,$$

$$(C12) \quad (x \rightarrow (y \rightarrow z)) \rightarrow ((x \rightarrow y) \rightarrow (x \rightarrow z)) ,$$

$$(C13) \quad ((x \rightarrow y) \rightarrow x) \rightarrow x ,$$

$$(R6) \quad \frac{x, x \rightarrow y}{y} \quad (\text{weak modus ponens}).$$

We denote by F_c and T_c the sets of all the formulas and all the formulas which can be obtain from R0,R1,C1,...,C5, respectively.

Let α, β be the formulas of F_c , we write $\alpha \equiv \beta$ if and only if $\alpha \rightarrow \beta, \beta \rightarrow \alpha \in T_c$.

Then, it holds ([10]) .

$$(R7) \quad x \equiv x ,$$

$$(R8) \quad \frac{x \equiv y}{y \equiv x} ,$$

$$(R9) \quad \frac{x \equiv y, y \equiv z}{x \equiv z} ,$$

$$(R10) \quad \frac{x \equiv y}{(x \rightarrow z) \equiv (y \rightarrow z)} ,$$

$$(R11) \quad \frac{x \equiv y}{(z \rightarrow x) \equiv (z \rightarrow y)} ,$$

In order to study the 3-ILPC with algebraic techniques, in 1968 A. Monteiro [10] introduced the notion of implicative three-valued Lukasiewicz algebra (or I₃-algebra) as algebras $(A, \rightarrow, 1)$ of type $(2,0)$ satisfying :

$$(I1) \quad 1 \rightarrow x = x ,$$

$$(I2) \quad x \rightarrow (y \rightarrow x) = 1 ,$$

$$(I3) \quad (x \rightarrow y) \rightarrow ((y \rightarrow z) \rightarrow (x \rightarrow z)) = 1 ,$$

$$(I4) \quad (x \rightarrow y) \rightarrow y = (y \rightarrow x) \rightarrow x ,$$

$$(I5) \quad ((x \rightarrow y) \rightarrow (y \rightarrow x)) \rightarrow (y \rightarrow x) = 1 ,$$

$$(I6) \quad ((x \rightarrow (x \rightarrow y)) \rightarrow x) \rightarrow x = 1 \quad (\text{See also } [2,3,4,5,6]).$$

We denote by I_s the variety of the I_s -algebras .

Next, we give the most simple and important example of an I_s -algebra.

Let $T = \{0, 1/2, 1\}$ and \rightarrow be the operation defined by means of the table :

\rightarrow	0	$1/2$	1
0	1	1	1
$1/2$	$1/2$	1	1
1	0	$1/2$	1

Then $(T, \rightarrow, 1) \in I_s$.

If $A \in I_s$, for all $x, y \in A$ we define the operation \rightarrow by means of the formula :

$$(I7) \quad x \rightarrow y = x \rightarrow (x \rightarrow y) \quad .$$

The following properties are valid in every I_s -algebra and they are proved in [10] :

$$(I8) \quad x \rightarrow x = 1 \quad ,$$

$$(I9) \quad x \rightarrow 1 = 1 \quad ,$$

(I10) The relation \leq defined by $x \leq y$ if and only if $x \rightarrow y = 1$ is a partial order on A .

$$(I11) \quad x \leq y \text{ implies } y \rightarrow z \leq x \rightarrow z \quad ,$$

(I12) (A, \leq) is a join semi-lattice and for all $x, y \in A$, the element $x \vee y = (x \rightarrow y) \rightarrow y$ is the supremum of x, y .

$$(I13) \quad x \rightarrow (y \rightarrow z) = y \rightarrow (x \rightarrow z) \quad ,$$

$$(I14) \quad (x \rightarrow y) \rightarrow ((z \rightarrow x) \rightarrow (z \rightarrow y)) = 1 \quad ,$$

(I15) $x \leq y$ implies $z \rightarrowtail x \leq z \rightarrowtail y$,

(I16) $x \rightarrow (y \rightarrow x) = 1$,

(I17) $(x \rightarrow (y \rightarrow z)) \rightarrow ((x \rightarrow y) \rightarrow (x \rightarrow z)) = 1$,

(I18) $((x \rightarrow y) \rightarrow x) \rightarrow x = 1$,

(I19) $1 \rightarrow x = x$.

If $A \in I_3$ and there exists an element $0 \in A$ such that

(I20) $0 \leq x$, for all $x \in A$, (i.e. A is bounded)

Then

(I21) $(A, \wedge, \vee, \sim, \nabla, 1)$ is a three-valued Lukasiewicz algebra where
 $\sim x = x \rightarrowtail 0$, $x \wedge y = \sim(\sim x \vee \sim y)$ [8,9,10] and it verifies
 $x \rightarrowtail y = (\nabla \sim x \vee y) \wedge (\nabla y \vee \sim x)$ ([10]) .

If $\mathcal{A} = (A, \rightarrowtail, 1) \in I_3$, we denote by $L(\mathcal{A})$ (or $L(A)$) the algebra

$(A, \wedge, \vee, \sim, \nabla, 1)$ described in (I21) .

Let $A \in I_3$. $D \subseteq A$ is a deductive system (d.s.) of A if it satisfies:

(D1) $1 \in D$,

(D2) $x, x \rightarrowtail y \in D$ imply $y \in D$.

Let $D(A)$ be the set of all d.s. of A .

By [10] (See also [4]) we know that (D2) is equivalent to:

(D'2) $x, x \rightarrow y \in D$ imply $y \in D$.

We denote by $Con_{I_3}(A)$ the set of all I_3 -congruences of A . If

$R \in Con_{I_3}(A)$, x_R represents the equivalence class of x , $x \in A$ and $q_R : A \rightarrow A/R$ defined by $q_R(x) = x_R$ is the canonical epimorphism.

Then :

$\text{Con}_{I_3}(A) = \{ R(D) : D \in D(A) \}$, where

$R(D) = \{(x,y) \in A^2 : x \rightarrowtail y, y \rightarrowtail x \in D\}$. Furthermore, for each $R \in \text{Con}_{I_3}(A)$ we have that $D = 1_R \in D(A)$ and $R = R(D)$.

Let $\text{Hom}_{I_3}(A,B)$ ($\text{Epi}_{I_3}(A,B)$) be the set of all the I_3 -homomorphisms (I_3 -epimorphisms) from A into B (from A onto B) .

If $h \in \text{Hom}_{I_3}(A,B)$ then the set $N(h) = \{ x \in A : h(x) = 1 \}$, called kernel of h , has the following properties :

(H1) $N(h) \in D(A)$,

(H2) $(x,y) \in R(N(h))$ if and only if $h(x) = h(y)$,

(H3) $A/N(h)$ and $h(A)$ are isomorphic I_3 -algebras

(i.e. $A/N(h) \cong h(A)$) .

We denote by $E(A)$ the set of all maximal d.s. (m.d.s.) of A .

Then it holds :

(M1) $\bigcap \{ M : M \in E(A) \} = \{1\}$,

(M2) For each $M \in E(A)$, $A/M \cong S$, where S is a non trivial I_3 -subalgebra of T .

Next , we consider the 3-SLPC which is an extension of the 3-ILPC because it is obtained by adding the connective \wedge and the axiom schemes :

(SC1) $(x \wedge y) \rightarrowtail x$,

(SC2) $(x \wedge y) \rightarrowtail y$,

(SC3) $x \rightarrowtail (y \rightarrowtail (x \wedge y))$,

(SC4) $((x \rightarrowtail y) \wedge (x \rightarrowtail z)) \rightarrowtail (x \rightarrowtail (z \wedge y))$.

We denote by F_s and T_s the sets of all the formulas of the 3-SLPC

and all the formulas which can be obtained from R0,R1,C1,...,C5,SC1,...,SC4, respectively .

If $\alpha, \beta \in F_s$ we write $\alpha \equiv_s \beta$ if $\alpha \rightarrowtail \beta, \beta \rightarrowtail \alpha \in T_s$.

Then , we have :

$$(SR1) \quad \frac{x \rightarrow y, x \rightarrow z}{x \rightarrow (z \wedge y)} : \quad$$

- (1) $(x \rightarrow y) \rightarrow ((x \rightarrow z) \rightarrow ((x \rightarrow y) \wedge (x \rightarrow z)))$ [SC3] ,
- (2) $(x \rightarrow y) \wedge (x \rightarrow z)$ [(1),hip,R1] ,
- (3) $((x \rightarrow y) \wedge (x \rightarrow z)) \rightarrow (x \rightarrow (z \wedge y))$ [SC4] ,
- (4) $x \rightarrow (z \wedge y)$ [(3),(2),R1] .

$$(SC5) \quad (x \rightarrow (y \wedge z)) \rightarrow ((x \rightarrow z) \wedge (x \rightarrow y)) : \quad$$

- (1) $(y \wedge z) \rightarrow z$ [SC2] ,
- (2) $(x \rightarrow (y \wedge z)) \rightarrow (x \rightarrow z)$ [(1),R4] ,
- (3) $(y \wedge z) \rightarrow y$ [SC1] ,
- (4) $(x \rightarrow (y \wedge z)) \rightarrow (x \rightarrow y)$ [(3),R4] ,
- (5) $(x \rightarrow (y \wedge z)) \rightarrow ((x \rightarrow z) \wedge (x \rightarrow y))$ [(2),(4),SR1] ,

$$(SC6) \quad (x \rightarrow y) \rightarrow ((x \wedge z) \rightarrow (y \wedge z)) : \quad$$

- (1) $(x \wedge z) \rightarrow x$ [SC1] ,
- (2) $(x \rightarrow y) \rightarrow ((x \wedge z) \rightarrow y)$ [(1),R3] ,
- (3) $(x \wedge z) \rightarrow z$ [SC2] ,
- (4) $(x \rightarrow y) \rightarrow ((x \wedge z) \rightarrow z)$ [(3),R2] ,
- (5) $(x \rightarrow y) \rightarrow (((x \wedge z) \rightarrow y) \wedge ((x \wedge z) \rightarrow z))$ [(2),(4),SR1] ,
- (6) $(((x \wedge z) \rightarrow y) \wedge ((x \wedge z) \rightarrow z)) \rightarrow ((x \wedge z) \rightarrow (y \wedge z))$ [SC4] ,

$$(7) ((x \rightarrow y) \rightarrow (((x \wedge z) \rightarrow ((x \wedge z) \rightarrow z))) \rightarrow ((x \rightarrow y) \rightarrow ((x \wedge z) \rightarrow (y \wedge z))) \quad [(6), R4] ,$$

$$(8) (x \rightarrow y) \rightarrow ((x \wedge z) \rightarrow (y \wedge z)) \quad [(7), (5), R1] .$$

$$(SC7) (x \wedge y) \rightarrow (y \wedge x) :$$

$$(1) (x \wedge y) \rightarrow x \quad [SC1] ,$$

$$(2) (x \wedge y) \rightarrow y \quad [SC2] ,$$

$$(3) (x \wedge y) \rightarrow (y \wedge x) \quad [(1), (2), SR1] .$$

$$(SC8) (x \rightarrow y) \rightarrow ((z \wedge x) \rightarrow (z \wedge y)) :$$

$$(1) (x \rightarrow y) \rightarrow ((x \wedge z) \rightarrow (y \wedge z)) \quad [SC6] ,$$

$$(2) (x \wedge z) \rightarrow ((x \rightarrow y) \rightarrow (y \wedge z)) \quad [(1), R5] ,$$

$$(3) ((z \wedge x) \rightarrow (x \wedge z)) \rightarrow ((z \wedge x) \rightarrow ((x \rightarrow y) \rightarrow (y \wedge z))) \quad [(2), R4] ,$$

$$(4) (z \wedge x) \rightarrow ((x \rightarrow y) \rightarrow (y \wedge z)) \quad [(3), SC7, R1] ,$$

$$(5) (x \rightarrow y) \rightarrow ((z \wedge x) \rightarrow (y \wedge z)) \quad [(4), R5] ,$$

$$(6) (y \wedge z) \rightarrow (z \wedge y) \quad [SC7] ,$$

$$(7) ((z \wedge x) \rightarrow (y \wedge z)) \rightarrow ((z \wedge x) \rightarrow (z \wedge y)) [(6), R4] ,$$

$$(8) ((x \rightarrow y) \rightarrow ((z \wedge x) \rightarrow (y \wedge z))) \rightarrow ((x \rightarrow y) \rightarrow ((z \wedge x) \rightarrow (z \wedge y))) \quad [(7), R4] ,$$

$$(9) (x \rightarrow y) \rightarrow ((z \wedge x) \rightarrow (z \wedge y)) \quad [(8), (5), R1] .$$

From the above results it follows at once that \equiv_s is a congruence on F_s and the Lindenbaum algebra $(F_s / \equiv_s, \rightarrow, \wedge, 1)$ (where $1 = T_s$) satisfies :

$$(1^*) (F_s / \equiv_s, \rightarrow, 1) \in Ia ,$$

(2°) The following identities are verified :

$$(S1) (\alpha \wedge \beta) \rightarrowtail \alpha = 1 ,$$

$$(S2) \alpha \rightarrowtail (\beta \wedge \gamma) = (\alpha \rightarrowtail \gamma) \wedge (\alpha \rightarrowtail \beta) .$$

2. SI₃-ALGEBRAS

2.1. DEFINITION. A SI₃-algebra is an algebra $(A, \rightarrowtail, \wedge, 1)$ (A, as short) of type $(2, 2, 0)$ where $(A, \rightarrowtail, 1) \in I_3$ and the following identities hold :

$$(S1) (x \wedge y) \rightarrowtail y = 1 ,$$

$$(S2) x \rightarrowtail (y \wedge z) = (x \rightarrowtail z) \wedge (x \rightarrowtail y) .$$

We denote by SI₃ the variety of the SI₃-algebras.

2.2. EXAMPLES.

(1°) Let $(T, \rightarrowtail, \wedge, 1)$, where $(T, \rightarrowtail, 1)$ is the I₃-algebra indicated in 1, and \wedge is defined by : $x \wedge y = 0$ for all $x, y \in T$. Then T verifies (S1) but not (S2), because :

$$1/2 \rightarrowtail (1 \wedge 1) = 1/2 \rightarrowtail 0 = 1/2 \text{ and } (1/2 \rightarrowtail 1) \wedge (1/2 \rightarrowtail 1) = 0.$$

(2°) Let $(T, \rightarrowtail, \wedge, 1)$, where $(T, \rightarrowtail, 1)$ is the I₃-algebra of 1, and \wedge is defined by the table :

\wedge	0	$1/2$	1
0	0	$1/2$	1
$1/2$	$1/2$	$1/2$	1
1	1	1	1

Then T satisfies (S2) but not (S1) because :

$$(1/2 \wedge 0) \rightarrow 0 = 1/2 \rightarrow 0 = 1/2 \neq 1$$

(3°) Let $T^* = (T, \rightarrow, \wedge, 1)$, where $(T, \rightarrow, 1)$ is the Is-algebra of 1,

and \wedge is defined by the table :

\wedge	0	$1/2$	1
0	0	0	0
$1/2$	0	$1/2$	$1/2$
1	0	$1/2$	1

Then T^* is an SI₃-algebra. The non trivial SI₃-subalgebras of T^* are $B^* = \{0, 1\}$ and $L^* = \{1/2, 1\}$. Moreover $B^* \cong L^*$.

2.3. LEMMA. For every $A \in \text{SI}_3$ it holds:

$$(S3) \quad 1 = 1 \wedge 1 ,$$

$$(S4) \quad y = 1 \wedge y ,$$

$$(S5) \quad x \wedge y = y \wedge x ,$$

$$(S6) \quad x \wedge y \leq x ,$$

$$(S7) \quad x \wedge y \leq y ,$$

$$(S8) \quad z \leq x , z \leq y \text{ imply } z \leq x \wedge y ,$$

$$(S9) \quad (A, \leq) \text{ is a meet semi-lattice where the infimum of } x, y \text{ is } x \wedge y ,$$

$$(S10) \quad (A, \vee, \wedge, 1), \text{ where } \vee \text{ is the operation determined by the formula (I12)}, \text{ is a distributive lattice with last element } 1 ,$$

$$(S11) \quad (x \rightarrow y) \rightarrow ((x \wedge z) \rightarrow (y \wedge z)) = 1 .$$

Proof.

- (S3) $1 = (I8) (1 \wedge 1) \rightarrow (1 \wedge 1) = (S2) ((1 \wedge 1) \rightarrow 1) \wedge ((1 \wedge 1) \rightarrow 1) = (S1) 1 \wedge 1$.
- (S4) $(1) y \rightarrow (1 \wedge y) = (S2) (y \rightarrow y) \wedge (y \rightarrow 1) = (I8, I9) = 1 \wedge 1 = (S3) 1$. From (1) and I10.
- (2) $y \leq 1 \wedge y$. By S1 and I10 (3) $1 \wedge y \leq y$. From (2) and (3) $y = 1 \wedge y$.
- (S5) $x \wedge y = (I1) 1 \rightarrow (x \wedge y) = (S2) (1 \rightarrow y) \wedge (1 \rightarrow x) = (I1) y \wedge x$.
- (S6) $1 = (S1) (y \wedge x) \rightarrow x = (S5) (x \wedge y) \rightarrow x$, then by I10 $x \wedge y \leq x$.
- (S7) It follows from (S1) and I10.
- (S8) Let $x, y, z \in A$ be such that (1) $z \rightarrow x = 1$, (2) $z \rightarrow y = 1$, then $1 = (S3) 1 \wedge 1 = ((1), (2)) (z \rightarrow x) \wedge (z \rightarrow y) = (S2, S6) z \rightarrow (x \wedge y)$, hence by I10 $z \leq x \wedge y$.
- (S9) It follows from (S6), (S7) and (S8).
- (S10) We shall prove that the cancelation law holds:
- (C.L.) $x \wedge y = x \wedge z$, $x \vee y = x \vee z$ imply $y = z$.
- Indeed, let $x, y, z \in A$ be such that:
- (1) $x \wedge y = x \wedge z$, (2) $x \vee y = x \vee z$, then
- (3) $x \rightarrow y = (x \rightarrow y) \wedge 1 = (I8) (x \rightarrow y) \wedge (x \rightarrow x) = (S2) x \rightarrow (x \wedge y) = ((1)) x \rightarrow (x \wedge z) = (S2, I8, S4) x \rightarrow z$.
- On the other hand, $1 = (I12, I10) y \rightarrow (x \vee y) = ((2)) y \rightarrow (x \vee z) = (I12) y \rightarrow ((x \rightarrow z) \rightarrow z) = (I13) (x \rightarrow z) \rightarrow (y \rightarrow z)$. By I10 $x \rightarrow z \leq y \rightarrow z$, so by (3) we have
- (4) $x \rightarrow y \leq y \rightarrow z$.

Furthermore $y \rightarrowtail x = (S4, S5)$ $(y \rightarrowtail x) \wedge 1 = (I8)$ $(y \rightarrowtail x) \wedge (x \rightarrowtail x) = (S2)$ $y \rightarrowtail (y \wedge x) = ((1), S5)$ $y \rightarrowtail (x \wedge z) = (S2)$ $(y \rightarrowtail z) \wedge (y \rightarrowtail x)$, and so (5) $y \rightarrowtail x \leq y \rightarrowtail z$. Then $1 = (I5, I12)$ $(x \rightarrowtail y) \vee (y \rightarrowtail x) = ((4), (5))$ $y \rightarrowtail z$, hence by I10 (6) $y \leq z$. In a similar way we prove (7) $z \leq y$.

From (6) and (7) $y = z$.

(S11) $(x \rightarrowtail y) \rightarrowtail ((x \wedge z) \rightarrowtail (y \wedge z)) = (S2)$ $(x \rightarrowtail y) \rightarrowtail (((x \wedge z) \rightarrowtail z) \wedge ((x \wedge z) \rightarrowtail y)) = (S1, S4)$ $(x \rightarrowtail y) \rightarrowtail ((x \wedge z) \rightarrowtail y) = (I13)$ $(x \wedge z) \rightarrowtail ((x \rightarrowtail y) \rightarrowtail y) = (I12)$ $(x \wedge z) \rightarrowtail (x \vee y) = 1$.

It is easy to see that :

2.4. LEMMA. Let $(A, \rightarrowtail, \wedge, 1) \in \text{SIs}$ and $(A, \rightarrowtail, 1)$ be bounded with first element 0. Then, it holds :

$$x \wedge y = \sim(\sim x \vee \sim y) = (((x \rightarrowtail 0) \rightarrowtail (y \rightarrowtail 0)) \rightarrowtail (y \rightarrowtail 0)) \rightarrowtail 0.$$

2.5. EXAMPLE. Let T^* be the SIs-algebra indicated in 2.2(3°), N be the set of the positive integers and T^{*N} be the set of all functions from N into T^* pointwise algebraized. For each $f \in T^{*N}$ we denote by A_f the set $\{ i \in N : f(i) \neq 1 \}$. Let $A = \{ f \in T^{*N} : |A_f| < \omega \}$. Then it is easy to see that A is a SIs-subalgebra of T^{*N} , so $A \in \text{SIs}$ and A does not have first element.

If $\mathcal{A} = (A, \rightarrowtail, \wedge, 1) \in \text{SIs}$ then we represent by \mathcal{A}^R the reduct $(A, \rightarrowtail, 1)$ and when there is no doubt we write A^R instead of \mathcal{A}^R .

2.6. LEMMA. If $A \in \text{SI}_3$ then $\text{Con}_{\text{SI}_3}(A) = \text{Con}_{\text{I}_3}(A^R)$.

Proof. It is clear that $\text{Con}_{\text{SI}_3}(A) \subseteq \text{Con}_{\text{I}_3}(A^R)$.

If $R \in \text{Con}_{\text{I}_3}(A^R)$ then there exists $D \in \mathbb{D}(A)$ such that $R = R(D)$. If $(x, y) \in R$ then (1) $x \rightarrowtail y \in D$, (2) $y \rightarrowtail x \in D$. From (1), (S11), D1, and D2 we have (3) $(x \wedge z) \rightarrowtail (y \wedge z) \in D$. In the same way we prove that (4) $(y \wedge z) \rightarrowtail (x \wedge z) \in D$. By (3) and (4) , (5) $(x \wedge z, y \wedge z) \in R$. By (5) and (S5) ,(6) $(z \wedge x, z \wedge y) \in R$. Hence, from (5) and (6) it results that $R \in \text{Con}_{\text{SI}_3}(A)$.

Therefore $\text{Con}_{\text{I}_3}(A^R) \subseteq \text{Con}_{\text{SI}_3}(A)$.

2.7. REMARKS. Since the SI_3 -congruences of a SI_3 -algebra A are determined by the d.s. of A and the operator \rightarrow verifies the properties (I16),...,(I19) then taking into account (D'2) and some results of A. Monteiro [11] (See also [12]) we can state that:

(1°) Every proper d.s. of A is the intersection of m.d.s. of A . In particular {1} is the intersection of all the m.d.s. of A .

(2°) In any SI_3 -algebra A the following conditions are equivalent

(a) A/M is a simple SI_3 -algebra,

(b) $M \in \mathbb{E}(A)$,

(c) If $a \in A-M$ and $b \in A$ then $a \rightarrow b \in M$.

(3°) The SI_3 -algebras are semi-simples in the sense of the following theorem.

2.8. THEOREM. Every non trivial SI₃-algebra A is a subdirect product of simple SI₃-algebras.

Proof.

Let $P = \prod_{M \in E(A)} A/M$ and $\varphi : A \rightarrow P$ be the mapping defined by

$\varphi(f) = F$, where $F(M) = q_M(f)$, for all $M \in E(A)$.

Since $q_M : A \rightarrow A/M$ is the canonical epimorphism, φ is a SI₃-homomorphism and by 2.7(1) φ is injective.

2.9. THEOREM. Let $A \in SI_3$ be non trivial. Then the following conditions are equivalent

(1°) A is simple,

(2°) $A \cong S$, where S is a non trivial SI₃-subalgebra of T^* .

Proof.

(1°) \Rightarrow (2°) If $A \in SI_3$ is simple, then it has more than one element. By 2.6 and the results indicated in 1 it follows that $D(A) = \{1\}$. Then A^R is a simple Is-algebra, so $A^R \cong S^R$, where S^R is an Is-subalgebra of T . Finally, taking into account that (A, \wedge, \vee) is a distributive lattice we obtain that $A \cong S$, where S is a non trivial SI₃-subalgebra of T^* .

(2°) \Rightarrow (1°) It is clear that T^* and all their non trivial SI₃-subalgebras are simple, so if A verifies (2°) then A is simple.

2.10. REMARK. If $A \in SI_3$ is finite then the application φ of 2.8

is onto. Indeed, since A is finite it has first element. Then we can consider the Lukasiewicz algebra $L(A)$ of (I21) and apply the known results for these algebras.

3. SI₃-ALGEBRAS WITH A FINITE SET OF FREE GENERATORS

Now, we are going to determine the structure of the SI₃-algebra with n free generators, n positive integer.

3.1. DEFINITION. If c is a positive cardinal number, we say that $S(c)$ is SI₃-algebra with c free generators if :

- (1) $S(c)$ has a set of generators G such that $|G| = c$,
- (2) Any function f from G into a SI₃-algebra A can be extended to a SI₃-homomorphism $h : S(c) \rightarrow A$.

Since the notion of SI₃-algebra is equationally definable, a result of G. Birkhoff [1] allows us to assert that for any cardinal $c > 0$ there exists $S(c)$ and it is unique up to isomorphisms. Moreover the SI₃-homomorphism of definition 3.1 is unique.

If we consider in the 3-SLPC a set of propositional variables of cardinal c , $c > 0$ then the Lindenbaum algebra mentioned in 1. is a SI₃-algebra with c free generators.

Let $G = \{g_1, g_2, \dots, g_n\}$ a set of free generators of $S(c)$. We shall denote by T^G the set of all functions from G into $T = \{0, 1/2, 1\}$.

Since each function $f \in T^G$ can be extended to a unique

$h \in \text{Hom}_{\text{SI}_3}(S(n), T^*)$ such that $h/G = f$, the correspondence $f \rightarrow h$ is bijection between T^G and $\text{Hom}_{\text{SI}_3}(S(n), T^*)$. Hence $\text{Hom}_{\text{SI}_3}(S(n), T^*)$ is finite.

3.2. LEMMA. The map $\varphi : \text{Hom}_{\text{SI}_3}(S(n), T^*) \rightarrow E(S(n))$ defined by $\varphi(h) = N(h)$ is onto .

Proof.

Let $M \in E(S(n))$ and $q_M : S(n) \rightarrow S(n)/M$ be the canonical epimorphism. From 2.7.(2°) and 2.9 there exists a SI_3 -homomorphism $i : S(n)/M \rightarrow T^*$. Then $h = i \circ q_M \in \text{Hom}_{\text{SI}_3}(S(n), T^*)$ and $\varphi(h) = M$, so φ is onto .

$S(n)$ is a subdirect product of finite algebras $S(n)/M$, $M \in E(S(n))$. Then, taking into account 3.2 and 2.9 we have :

3.3. COROLLARY. If n is a positive integer, then $S(n)$ is finite and $S(n) \cong \prod_{M \in E(S(n))} S(n)/M$.

Let $E_1 = \{ M \in E(S(n)) : |S(n)/M| = 2 \}$ and

$E_2 = \{ M \in E(S(n)) : |S(n)/M| = 3 \}$. Then $\{E_1, E_2\}$ is a partition of E . Taking into account 2.2(3°), 2.7(2°) and 2.9 we can write :

$$S(n) \cong \prod_{E_1}^{|\mathbb{E}_1|} \times \prod_{E_2}^{|\mathbb{E}_2|} \quad (I).$$

Now, we prove :

3.4. LEMMA. For each non trivial $h \in \text{Hom}_{\text{SI}_3}(S(n), T^*)$ there exists $h' \in \text{Hom}_{\text{SI}_3}(S(n), T^*)$ which verifies :

- (1°) $h'(S(n)) = \{1/2, 1\} = L^*$,
- (2°) $N(h') = N(h)$.

Proof. Let $h \in \text{Hom}_{\text{SI}_3}(S(n), T^*)$ be non trivial and suppose that $h(S(n)) = \{0, 1\} = B$.

Then defining $\alpha : B \rightarrow L$ in the following way : $\alpha(0) = 1/2$, $\alpha(1) = 1$ we have that $h' = \alpha \circ h$ verifies (1°) and (2°) .

Let $A = \{ h \in \text{Hom}_{\text{SI}_3}(S(n), T^*) : h(S(n)) = L^* \}$, $B = \text{Epi}_{\text{SI}_3}(S(n), T^*)$

Since T^* has non trivial automorphism, from 3.4. it results :

$$|E_1| = |A| \quad (\text{II}) \quad ,$$

$$|E_2| = |B| \quad (\text{III}).$$

On the other hand , let $A' = \{ f \in T^G : (f(G)) = L^* \}$ and $B' = \{ f \in T^G : (f(G)) = T^* \}$, where $(f(G))$ is the SI_3 -subalgebra of T generated by $f(G)$. Then we have :

$$|A| = |A'| \quad (\text{IV}) \quad ,$$

$$|B| = |B'| \quad (\text{V}) \quad .$$

Furthermore,

$$A' = \{ f \in T^G : f(G) \subseteq \{1/2, 1\} \subseteq T \text{ y } 1/2 \in f(G) \}.$$

Let $B'_1 = \{ f \in B' : f(G) = \{0, 1/2\} \}$ and

$B'_2 = \{ f \in B' : f(G) = \{0, 1/2, 1\} \}$. Then $\{B'_1, B'_2\}$ is a partition of B' .

Finally we have

$$|A'| = 2^n - 1 \quad (\text{VI}) \quad ,$$

$$|B'_1| = 2^n - 1 ,$$

$$|B'_2| = \sum_{i=0}^{3-1} (-1)^i \binom{3}{i} (3-i)^n = 3^n - 3 \cdot 2^n + 3 ,$$

$$|B'| = (2^n - 1) + (3^n - 3 \cdot 2^n + 3) = 3^n - 2^{n+1} + 2 \quad (\text{VII})$$

From (I), (II), ..., (VII) we obtain :

3.6. THEOREM. The SI₃-algebra S(n) with n free generators, n integer, n > 0 verifies :

$$(1^\circ) \quad S(n) \cong L^{\star 2^{n-1}} \times T^{\star 3^n - 2^{n+1} + 2} ,$$

$$(2^\circ) \quad |S(n)| = 2^{2^{n-1} \cdot 3^{n-1} - 2^{n+1} + 2} .$$

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